

Digital Currency Generation Using Hash Function and Digital Signature (DigCur)

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Abstract

This paper proposed a strategy to create a digital currency by producing a list of serial numbers by an authorized organization (Issuer). Each number in the serial numbers list will be merged with the issuer's secret key (Sec_key) and then hashed using hash function to produce the issuer hash code (H-code). The issuer hash code will be considered as a digital banknote and then it will be distributed to the payee as a type money. When a person likes to pay (transfer) these digital banknotes to another person, he/she must prove this transaction by his/her private key using digital signature process. This mechanism needs a third party to control the process of money transaction by checking the validity of the digital signature and who is authorized to own this digital banknote. The third party is the issuer himself (the central bank). For each digital banknote the issuer hash code will be saved in a record without changing to check if this digital banknote is counterfeit or not. The problem of double spending will be solved by swapping the owner identifier immediately after the transaction process. The proposed method aims to produce a digital cryptocurrency to be an alternative to the traditional physical currencies and it will be called as DigCur. According to some tests the proposed method can be considered a good method to take advantage of it instead of other cash applications. [DOI: [10.22401/ANJS.22.2.09](https://doi.org/10.22401/ANJS.22.2.09)]

Keywords: Digital Currency, Secret Key, Hash Function, Digital Signature, Digital Banknotes.

1- Introduction

Digital currency is a type of money but it is formed digitally, also it is a payment method represented in digital form [3]. The digital currency has become a reality, and the modern world requires faster and safer financial transactions, it is therefore necessary to deal with mechanisms which make money transaction more efficient, and this paper proposed a method for generating digital currency with some principles of cryptography.

There are many ways to transfer money from sender to the receiver and all of them depend on the amount of money that will transfer but not the banknote itself, for example; in any electronic payment system to transfer 200\$ from the sender's wallet to the receiver's wallet the process required a subtraction and addition operations with identification and authentication processes, so all electronic financial transactions not take the money units (banknotes) in the transaction process. The method suggested in this research deal with the banknotes as digital currency and will be called as DigCur.

2- Digital Currency

There are two terms; virtual currency and digital currency, virtual is anything created from nothing, while digital is a process of generating the currency by the computer, thus the term digital is preferred over virtual. In short, the two terms refer to the alternative method of conventional currency [3]. Virtual currencies issued by private agents and there is no control by the governments [15], therefore DigCur is a type of digital currency. In this research the digital currency that will be generated depends on the serial number which will be determined by the central authority (Central Bank).

There are many cryptocurrencies and bitcoin is one of them. Bitcoin depends on special methodology in which; the electronic cash for online payments send the money directly from one point to another without going through the third party [8], also the bitcoin paper proposed a strategy for solving the problem of double spending. Bitcoin is good and new method for cryptocurrency but it is not safe enough to be replacement for real money [10].

Centralized payment systems solve the problem of double spending but it needs some trust by the customers [2], for this reason we need the third party to control the works.

There are some differences between cryptocurrency and digital currency. Cryptocurrency is digital currency in its essence but it is differ in the *Structure, Anonymity, Transparency and Transaction manipulation*. The digital currency is centralized while cryptocurrency is decentralized, identification is required in digital currency, and in cryptocurrency identification is not required. The process of money transaction in digital currency is not transparent it gives some privacy for the customer wallet, in cryptocurrency the transaction can be seen by anyone because of the public chain. In the digital currency there is a full control by the third party, in cryptocurrency there is no control on the transaction [13].

One of the first attempts of electronic cash is explained in [6] where the researchers applied the digital signature to verify the coins, the system works in offline and online with a centralized database to solve the problem of double spending, the system can detect the double spend in online manner only, but in [6] there is no clear strategy for digital currency generation.

3- The purpose of generating DigCur

Everything in this life is in developing, banking and finance is one of the most applications that must keep up this development. *This paper aims to generate a digital currency using the principles of the cryptography*. The motives for proposing a new strategy (DigCur) for generating digital currency which depends on hash code and digital signature are:-

- 1- Safety of currency fraud (no chance for counterfeit money).
- 2- There is no dual digital banknote (uniqueness).
- 3- Digital signature prevents the sender from denying the process of money transaction that accomplished by him.
- 4- Checking the sender and receiver identities (identification and authentication).

- 5- Control the financial transfer.
- 6- The third party prevents the double spending.
- 7- Ease of dealing with banking operations.
- 8- No need for cash money.
- 9- New strategy for financial operations.

4- DigCur Mechanism

DigCur is digital currency because it is generated by computer program, and it is Cryptocurrency because it is needs to produce a hash code and digital signature. Sections in below explains how the DigCur will be generated and how the persons can transfer DigCur to others.

4-1 DigCur generation by the central bank

Central bank is the only authorized committee that must be able to generate DigCur, the process of DigCur generation will begin by selecting an initial value to produce a list of serial numbers, each serial number (SN) will be converted to a string and then it will be merged with the secret key (*Sec_key*) that has been selected by the central bank. The hash function will proceed on the merged string to generate the hash code, $H_code = HF("SN" + "Sec_key")$, where *HF* is the hash function, *SN* is the serial number and *Sec_key* is the secret key chosen by the issuer, in this research MD5 is the method that has been selected to generate the hash code because of its simplicity. MD5 is faster than SHA [14]. MD5 is described in [11]. In Microsoft Visual basic.net 2013 there is a built in function make easy to compute the MD5 of any input, we used it to compute the hash code.

At this stage the *H_code* is the core. This manner will simplify the operation of checking whether this currency is counterfeit or not, the process of detecting currency fraud is usually done by a third party (the central bank) which owns the secret key (*Sec_key*), where the central bank is the issuer himself.

To explain how the DigCur will be generated:- the issuer must select a start point and end point to generate a list of serial numbers and then each one of the serial numbers will be manipulated to generate the DigCur.

For example let;

SN="44670"

Sec_key="hello"

MergeSN_Sec_key="44670hello"

Then;

H_code=MD5(MergeSN_Sec_key)

H_code=76da990ed294983ef68c794243afb50

8 → hash code of 32 Hexadecimal digits, this hash code is unique. And so on for the next serial number (SN).

Table (1) shows a list of digital banknotes that has been generated by the issuer program.

Each digital banknote is DigCur, and there are many classes of DigCur.

All these digital banknotes are ready to be distributed to the payees, where each row in the table above is a record for one digital banknote.

It is good to insert a symbol with the serial numbers to give a uniqueness to these numbers and this may be done by agreement all the countries about these symbols, as in the case of phone country code or country internet codes (for example; IQ44677), IQ is a symbol for Iraq and 44677 is the SN.

Table (1)
Some digital banknotes of class 5DigCur with secret key ("hello").

SN	Date of Issue	Currency class	Issuer	H_code
44670	04/06/2018	5DigCur	Z_BANK	76da990ed294983ef68c794243afb508
44671	04/06/2018	5DigCur	Z_BANK	aefed49a7eb272d8f0ac533045fef455
44672	04/06/2018	5DigCur	Z_BANK	50004773951ec9817b03771f5995b228
44673	04/06/2018	5DigCur	Z_BANK	e748199d84786663e452aacd1de40520
44674	04/06/2018	5DigCur	Z_BANK	92318e9299d206e15b068ea264805bb0
44675	04/06/2018	5DigCur	Z_BANK	682e44d04d0a76f35dbfd8c9dfb3d05b
44676	04/06/2018	5DigCur	Z_BANK	e5d344f51740e6ea518a9452ade43e23
44677	04/06/2018	5DigCur	Z_BANK	30159aa4c6c5060528fdbc26e0cca48
44678	04/06/2018	5DigCur	Z_BANK	17a73eb239a5e39e9e2003cff3b8e0d8
44679	04/06/2018	5DigCur	Z_BANK	d3a5d13b9f1404e865a1541d3e1d7f0f

4-2 DigCur distribution and transaction

After generating the digital banknotes of the DigCur, the list of these digital banknotes will be saved and then will be ready to the process of distribution by the central bank to the payees after signing each digital banknote. In this research work the Elgamal digital signature is the method that has been selected to sign the DigCur, therefore each template of the DigCur will contain the digital signature of the sender and other required data as shown in Table (2). The Elgamal digital signature algorithm is explained in [9][4][7]. In each transaction the process of digital signature will be done by the sender to validate the process of the money transaction. The digital signature will prevent the sender to deny the process of money transfer. Also the DigCur template must have the owner identifier (Owner ID) and this field is dynamic (exchanged with each DigCur transaction).

Table (2)
DigCur template.

Serial number	Date of Issue	Currency Class	Issuer	H_code	Digital signature	Owner ID
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4-3 Digital Signature and Hash Code.

Digital signature is the process of signing a message with sender's private key. RSA, Elgamal are some examples of digital signature methods that originally are public key encryption methods [9][7]. Digital signature is a process to protect the data from forgery and to prove the integrity of the data. Digital signature is a form of cryptographic transaction that makes the receiver to know the source of the message [12].

Hash function is an algorithm that can summarize the message of variable length to a unique and fixed size of *n* bits, MD5 will produce a code of 128 bits, SHA1 will produce a code of 160 bits [5]. Hash function is one way encryption, for a message *x*, and it is impossible to find *x* by reversing the parameters [12], and this will protect the secret

key (*Sec_key*) that must be known for the third party only.

Hash function is used in digital signatures and in data integrity [5][1], to simplify the process of digital signature the message must be hashed to produce the hash code, and this because it is useless to sign each character in the message, therefore no need to signing the entire message, but only the hash code of the entire message will be signed, and then the hash code with the parameters of the digital signature will be sent to the receiver, and the validity checking will be done by the receiver, but in this research work the digital signature validation will be done by the central bank (issuer) who is considered as a third party.

MD5, SHA-1, SHA-2 are some examples of message hashing functions [9][1]. MD5 and SHA-0 both of which have been broken [12], therefore we advise to use another hash function such as SHA3.

To compute the hash_code it must to apply the following formula:-

$$\text{hash_code} = \text{HF}(x)$$

As shown in section 4-2 the sender must signs each digital banknote, the message that must be signed is the *H_code* itself. The role of the hash function in this research is to verify the integrity of the digital banknotes (DigCur) and to apply the digital signature on the hash code of the message instead of an entire message, where the message (*x*) is the *SN* and the *Sec_key*.

4-4 Solving the problem of double spending

Double spending is the problem of using the same digital currency more than once [16]. In this paper, to solve the problem of double spending; the unique identifier (ID) of the person that will receive the digital banknote will be inserted in DigCur template instead of the sender's ID, for this the central bank will be responsible for matching the saved template (DigCur record) with the current transferred DigCur to decide if this person is authorized for transaction or not. To make the DigCur transaction more realistic we need a central database that will save copies of the DigCurs for all persons. In the central database there is a hierarchy for indexing the DigCurs based on the serial number to make the DigCur record is

easy for retrieving. For each DigCur record and after finishing the transaction process the sender's ID will be swapped immediately by the receiver's ID, and this will deny the sender to send the DigCur of a specific serial number again because of the matching and authorization process will be done by the central bank. Fig.(1) shows the mechanism of the DigCur transferring.

5- The proposed DigCur Algorithm

Sections below show the steps required for each stage in the DigCur.

5-1 DigCur generation algorithm

Steps below show the algorithm of generating the DigCur by the central bank (issuer).

Algorithm (1): DigCur Generation Algorithm

Input: Start_SN, End_SN, Sec_Key, DigCur class.

Output: DigCur records.

Step 1: Begin

Step 2: Set the range of the serial numbers between (Start_SN) and (End_SN).

Step3: Select the secret key (Sec_Key).

Step4: Set the DigCur class.

Step5: For each number (SN) in the range between (Start_SN) and (End_SN) must be converted to a string (st).

Step6: Merge (st) with (Sec_Key) to create the message (M).

Step7: Compute the hash code (H_code) by applying the hash function (HF) on the message (M).

Step8: Insert the (H_code) in the DigCur record (template).

Step9: For the next (SN) go to step 5.

Step10: Save the DigCur records, where each record must contains some necessary fields (SN, Date of Issuing, Currency Class, Issuer, H_code, empty cells for the next digital signature, and Owner ID).

Step11: End DigCur generation algorithm.

5-2 DigCur transaction algorithm

Steps below show the algorithm of DigCur transaction operation between the sender and the receiver.

Algorithm (2):

DigCur Transaction Algorithm.

Input: *SenPrivate_Key*, *SenPub_Key*, *DigCur* class, *Number of DigCur*.

Output: *Accept or reject the DigCur*.

Step1: Begin

Step2: Set the private key (*SenPrivate_Key*) and the public key (*SenPub_Key*) by the sender for digital signature operation.

Step3: Set the class and the number of the *DigCur* by the sender.

Step4: Prepare the *DigCur_i* for transmission.

Step5: Apply the digital signature (DS) on (*H_code*) using (*SenPrivate_Key*) to compute the (*V*).

Step6: Send the *DigCur* to the receiver (as a record) with (*V*).

Step7: The third party will merge (*SN*) with the (*Sec_Key*) and then compute the hash function (*HF*).

Step8: If (*HF*) is equal to (*H_code*) then the *DigCur* is valid, otherwise reject the *DigCur* transaction (*DigCur* is forged) and go to step 14.

Step9: Checking the digital signature validation of the sender by computing the (*Value1* and *Value2*) using (*SenPub_Key*).

Step10: If (*Value1*) is equal to (*Value2*) then the digital signature is valid and the transaction must be accomplished, otherwise reject the *DigCur* and go to step 14.

Step11: If the sender's ID match with the copy of the *DigCur* in the central database then immediately updates the *DigCur* record by assigning the receiver's ID in the owner ID cell, otherwise reject the transaction (unauthorized sender).

Step12: Update the sender and receiver wallets by the third party.

Step13: For the next *DigCur* go to step 4.

Step14: End *DigCur* transaction.

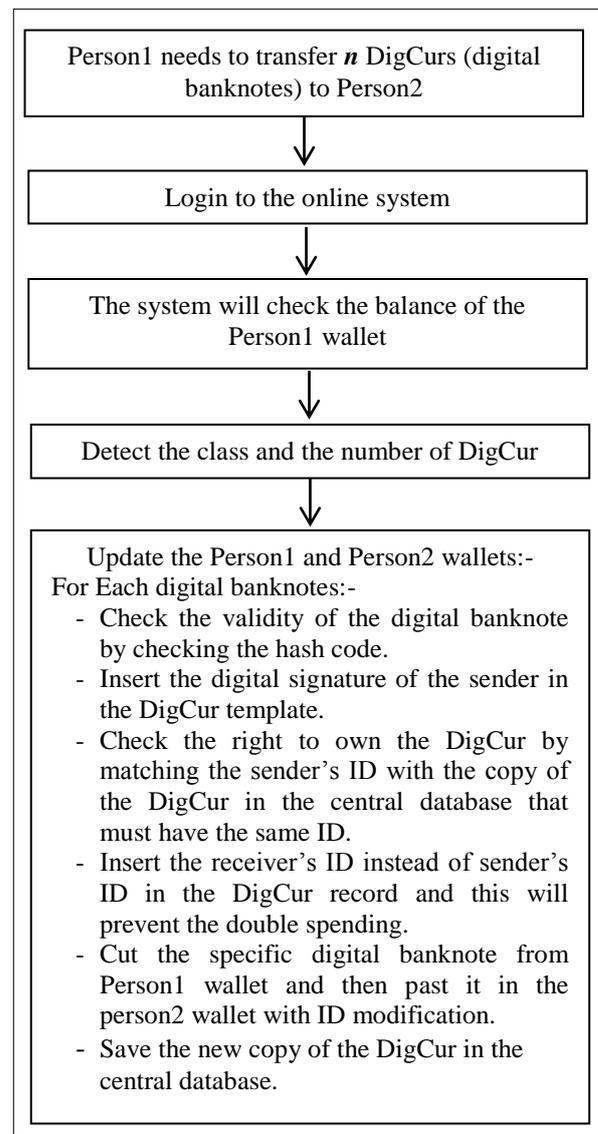


Fig.(1): The mechanism of the DigCur transferring.

6- Implementation Results

The proposed algorithm (DigCur algorithm) was implemented using Microsoft VB.Net2013, sections bellow show the results of tracing the system.

6-1 Generating DigCur

As shown is sections above, the central bank (CB) is the only authorized foundation that able to issue the *DigCur*, suppose the CB likes to issue 20 digital banknotes of class 50 *DigCur* with a secret key “system2016YaQu”, Table (3) shows the record of each digital banknote and then each record will be transacted to the payee. In the stage of *DigCur* generating there is no need to the digital signature, digital signature will be necessary in the transaction step.

Table (4) shows the *DigCur* with the secret key “system2016YaQ3”, you notice there are

new list of hash code differ totally from Table (3), and this because of hash function specification (MD5 in this case).

After signing each digital banknote, the central bank is ready to distribute the DigCur(s) to the payees.

Table (3)
List of DigCur of a proposed class (50DigCur) with secret key “system2016YaQu”.

SN	Date of Issue	Currency class	Issuer	H_code	Owner ID
432431	14/4/2018	50DigCur	ICB	89bf75b94e3f98bd713a64efb74fd165	ICB
432432	14/4/2018	50DigCur	ICB	6b36d13bbfcb86d924f073a54acae056	ICB
432433	14/4/2018	50DigCur	ICB	af9ce04bed31fcedbdf1f13d0eb4964b	ICB
432434	14/4/2018	50DigCur	ICB	e03ec1d88066af68cd00fecbe3589ab	ICB
432435	14/4/2018	50DigCur	ICB	e24d575e06ee83ce9f6ce3a72d203406	ICB
432436	14/4/2018	50DigCur	ICB	f19e24b78376b0fd6607fea7d9940c13	ICB
432437	14/4/2018	50DigCur	ICB	9928b8d5c72a84241fbb695cc47b95af	ICB
432438	14/4/2018	50DigCur	ICB	5fa576f05f5016f6d56988cce458b0f0	ICB
432439	14/4/2018	50DigCur	ICB	fb8c72749ac2d633ca9697e7cfd65a47	ICB
432440	14/4/2018	50DigCur	ICB	e2570555a22c830db807b3e03ba9d884	ICB
432441	14/4/2018	50DigCur	ICB	bdfd7b4d355c851a23ac9d7d3387a886	ICB
432442	14/4/2018	50DigCur	ICB	228d3b5e760f4c3c831080c73bca1612	ICB
432443	14/4/2018	50DigCur	ICB	47990c8596a743d54b040b13dfe02bb7	ICB
432444	14/4/2018	50DigCur	ICB	deafabb739ff9024d2087df4d44f46b4	ICB
432445	14/4/2018	50DigCur	ICB	1ac19b991e0c8c61e39470540976df1c	ICB
432446	14/4/2018	50DigCur	ICB	b2fac2f37643a9b2d6812b81d70afbdf	ICB
432447	14/4/2018	50DigCur	ICB	d6259f2bf1246b9207de901382ea7f4c	ICB
432448	14/4/2018	50DigCur	ICB	4accb387233d85cd4169f661a030d6c5	ICB
432449	14/4/2018	50DigCur	ICB	6a18fe867043dbf209a0313e76684120	ICB
432450	14/4/2018	50DigCur	ICB	4b5e81ab1d0d1135d04ad85d1a91db14	ICB

Table (4)
New list of DigCur of a proposed class (50DigCur) with secret key “system2016YaQ3”.

SN	Date of Issue	Currency class	Issuer	H_code	Owner ID
432431	14/4/2018	50DigCur	ICB	c86469d26ea89b3c0ae1231ca1109f14	ICB
432432	14/4/2018	50DigCur	ICB	710200e37d9909fdbc467bf41a0fc397	ICB
432433	14/4/2018	50DigCur	ICB	b4740aeb19f529820ec03d1f6293899f	ICB
432434	14/4/2018	50DigCur	ICB	c2cb69005ced89111e2cc0b9a8053f09	ICB
432435	14/4/2018	50DigCur	ICB	c3b0bfe21e7797c9e3cbfe092898d8cd	ICB
432436	14/4/2018	50DigCur	ICB	87ebc7e1a956d070de10601e4cf5d2f6	ICB
432437	14/4/2018	50DigCur	ICB	d4717029f87c0bb798e9efff8961aa0f	ICB
432438	14/4/2018	50DigCur	ICB	0888f559bf81afb368a3d1e2f4860248	ICB
432439	14/4/2018	50DigCur	ICB	d25fa673504c9f4def884869ee1667dc	ICB
432440	14/4/2018	50DigCur	ICB	46d87ac58a755e11f151d50ce0991e4c	ICB
432441	14/4/2018	50DigCur	ICB	03b168b068a97b8ba40cf12124be665c	ICB
432442	14/4/2018	50DigCur	ICB	9e23a552b6a2896875d2d3e296b7b7ae	ICB
432443	14/4/2018	50DigCur	ICB	a5cbe177025615b9b29e4e17a924bb5c	ICB
432444	14/4/2018	50DigCur	ICB	a29b34a576ddf670eb50e7796995dd06	ICB
432445	14/4/2018	50DigCur	ICB	d2d93c40a6f07a9cfda206c17ff62776	ICB
432446	14/4/2018	50DigCur	ICB	8cf047150527cd7f50717a5356019eab	ICB
432447	14/4/2018	50DigCur	ICB	4d82db92465b527a516a3b33b87d08d9	ICB
432448	14/4/2018	50DigCur	ICB	362ddc5e108597e6f1a2bb50eab0a327	ICB
432449	14/4/2018	50DigCur	ICB	ad3ffdf8069c2761ca0998b492505e19	ICB
432450	14/4/2018	50DigCur	ICB	0e3d85a201c385f663afbebc6493f187	ICB

6-2 DigCur transaction

To explain how the DigCur will transfer from one person to another, suppose the person (Person1) owns the following digital banknote and he wants to transmit it to (Person 2).

SN	432432
Date of Issue	14/4/2018
Currency class	50DigCur
Issuer	ICB
H_code	710200e37d9909fdb467bf41a0fc397
Owner ID	Person1

First the sender must sign each digital banknote and then the third party will complete the mechanism of the money transaction.

In this paper we select Elgamal digital signature method because of its simplicity.

To apply the Elgamal digital signature there are some computations (in below a tracing of the Elgamal Digital signature):-

In the sender side:-

Let $p=1753$ (p is primary number, to increase the complexity; p must be large prime number)

$a=8$ (a is primitive element to Z_p)

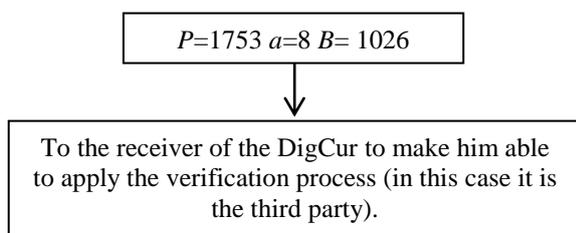
$d=12$ (d must be between (2) an ($p-2$)), it is the sender's private key.

Compute:-

$$B=(a^d \text{ mod } p) = 1026$$

(P , a , and B) is Peron1's public key, where Person1 is the sender.

Send p , a , and B as public key to the receiver to verify the digital signature.



Now it must to compute the digital signature of the message, the message (x) is the hash code of the digital banknote.

For the digital banknote in the above;

$x=710200$ $e37d99$ $09fdb4$ $467bf4$ $1a0fc3$ 97 and to simplify the computations and to overcome the problem of very large integer

numbers with exponential; we partitioned the message (x) to six sub-messages ($x_1, x_2, x_3, x_4, x_5,$ and x_6), first five sub-messages of size 24 bits, the last sub-message of size 8 bits, the total number of bits is 128 bits (MD5 hash code size).

$$x_1=(710200)_H = (7406080)_{10}$$

$$x_2=(e37d99)_H = (14908825)_{10}$$

$$x_3=(09fdb4)_H = (654780)_{10}$$

$$x_4=(467bf4)_H = (4619252)_{10}$$

$$x_5=(1a0fc3)_H = (1707971)_{10}$$

$$x_6=(97)_H = (151)_{10}$$

The third party will apply the two steps of the verification:-

First; it will check the hash code if it is accepted or not, by inserting the central bank secret key (Sec_key) "system2016YaQ3" with the SN "432432" and then compute the hash code, if it is accepted, then the second process will begin by checking the digital signature using the sender's public key (1753 8 1026).

Now; Let $Ke=7$, where $\text{gcd}(Ke, p-1)=1$

For the sub-message ($x_1=7406080$):-

Compute r and s (in the sender's side);

$$r=a^{Ke} \text{ MOD } P$$

$$r=8^7 \text{ MOD } 1753$$

$$r=2097152 \text{ MOD } 1753 = 564$$

$$s=((x-d*r)*Ke^{-1}) \text{ MOD } (p-1)$$

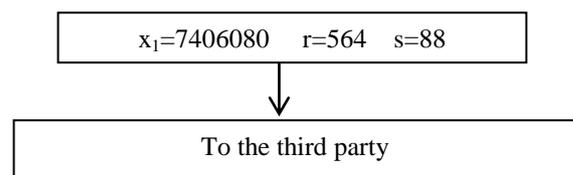
Ke^{-1} is the inverse of the Ke , $Ke^{-1}=751$

$$s=((7406080-12*564)*751) \text{ MOD } (p-1)$$

$$s=(7406080-6768)*751 \text{ MOD } 1752$$

$$s=5556883312 \text{ MOD } 1752=88$$

With each sub-message (x_i) there is a new s value, but r is the same.



The third party will check the validity by the verification process as follow:-

The third party will compute T_g and T_y ;

$$T_g=(B^r * r^s) \text{ MOD } p$$

$$T_g=(1026^{564} * 564^{88}) \text{ MOD } 1753=310$$

$$T_y=a^x \text{ MOD } p$$

$$T_y=8^{7406080} \text{ MOD } 1753=310$$

Since $T_g=T_y$ then the digital signature of the x_1 is valid.

And so on for the $x_2, x_3, x_4, x_5,$ and $x_6,$ the computations in the above will be repeated six times for each DigCur with the same p, a, d, B, ke and ke^{-1} .

In our system the template of the DigCur will be as follow:-

Serial number
Date of Issue
Currency class
Issuer
H_code (x):- ($x_1, x_2, x_3, x_4, x_5, x_6$)
Digital signature:- (p, a, B) ($r, s_1, s_2, s_3, s_4, s_5, s_6$)
Owner ID

We used the mechanism of partitioning because of the problem of very large integer that our system can't solve this problem.

For only one DigCur of class (50DigCur) will be as follow:-

432432
14/4/2018
50DigCur
ICB
710200e37d9909fdbc467bf41a0fc397
Digital Signature of the person (Person1) is:- $P=1753 a=8 B= 1026 r=564$ $s_1=88 s_2=1687 s_3=468 s_4=1676 s_5=1253$ $s_6=1057$
Person1

For each s_i the third party must compute T_g and T_y , and so on for all the digital banknotes (DigCurs) that will be transmitted from Person1 to Person 2. In reality the main role of the third party (the central bank) is to check the hash code (MD5) because it only knows the secret key by merging the serial number with the secret key and then compute the hash code. The process of digital signature it is possible to be in the receiver side, but it is good to be accomplished in the third party's servers (high speed computer). For each digital banknote the Owner ID will be updated to be

as a receiver ID immediately after checking the hash code and the digital signature, in this case it will be Person 2, also there is a need for updating the copy of this digital banknote in the central database to make the IDs matching in the next transaction is possible.

For testing; suppose there are four DigCurs sent by per1 to per2 as shown below:-

First DigCur

233432
21/11/2018
1DigCur
ICB
c10cc4df0ff957c2964c0a102713ad88
Digital Signature of the person (Person1) is:- $P=1753 a=8 B= 1497 r=564$ $s_1=152 s_2=67 s_3=1526 s_4=1668 s_5=759$ $s_6=1348$
Per1

Second DigCur

877613
21/11/2018
1DigCur
ICB
a9588d0e6814577937fe4b73b0e460c8
Digital Signature of the person (Person1) is:- $P=1753 a=8 B= 1497 r=564$ $s_1=983 s_2=552 s_3=845 s_4=1641 s_5=1052$ $s_6=356$
Per1

Third DigCur

655454
21/11/2018
1DigCur
ICB
fc09e4486aee5cf64133d7f13708d4f9
Digital Signature of the person (Person1) is:- $P=1753 a=8 B=1497 r=564$ $s_1=936 s_2=1694 s_3=1315 s_4=291 s_5=632$ $s_6=507$
Per1

Fourth DigCur

003343
21/11/2018
1DigCur
ICB
84315070b031724e41e212a845dc577e
Digital Signature of the person (Person1) is:- $P=1753$ $a=8$ $B=1497$ $r=564$ $s_1=540$ $s_2=643$ $s_3=1019$ $s_4=220$ $s_5=637$ $s_6=846$
Per1

The secret key of the central bank (*Sec_key*) in all cases must be “test1978”, the sender’s private key is ($d=27$), $Ke=7$, and the sender’s public key is ($p=1753$, $a=8$, $B=1497$).

The first DigCur will be accepted because the *H_code* is valid and the digital signature for each s_i is valid also (T_g equal to T_y).

The second DigCur will be rejected because the *H_code* is not valid although the digital signature is valid.

The third DigCur will be rejected because the digital signature is not valid (unauthorized person) although the *H-code* is valid.

The fourth DigCur will be rejected because the system when checked it in the central database did not find this DigCur subordinate to Per1 (ID matching).

7- Conclusions

This paper aims to propose a method for generating a digital currency that is useful for daily financial transactions and make it as an alternative of the traditional currencies. Also this paper may help the central banks to issue their own digital currency. The proposed digital currency (DigCur) is digital currency not virtual currency because it is generated by the computer program with a full control by the central bank, whereas the central bank is the third party that will be responsible for verifying the process of the digital money transaction. The digital moneys are units called as digital banknotes. This work generates a digital currency that is difficult to counterfeit as in the case of *paper* banknotes. The problem of double spending is solved by the central bank by checking the user identifier and makes the required update directly after finishing the transaction process. DigCur gives a control on the money transactions by using a

centralized database. Any centralized digital currency such as DigCur will reduce the concern from decentralized currencies, and this will protect the authority of the central banks. The hash code that is generated by a hash function is very helpful because of its uniqueness, whereas for each digital banknote (DigCur) there is a specific hash code. Also in this work we apply the digital signature to prevent the sender from denying the process of DigCur transfer. Also with some tests is possible to conclude the proposed method is good and novel idea for generating a centralized digital currency.

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